

# Information Leakage in Encrypted Network Traffic

#### **Attacks and Countermeasures**

#### Scott Coull RedJack

Joint work with:

- Charles Wright (MIT LL)
- Lucas Ballard (Google)
- Fabian Monrose (UNC)
- Gerald Masson (JHU)



#### **Encrypted Network Traffic**



#### **<u>Problem</u>: How do we stop eavesdropping?**

#### **Cryptography to the rescue!**





# **Encrypted Network Traffic**



- Encrypt the payload contents
- What about features of the network traffic itself?
  - Port numbers, packet sizes, timing info



- Traffic features may introduce a communications channel
  - Like a side-channel attack or covert channel





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# **REDJACK Preventing Information Leakage**

- Information theory says we should make the distributions uniform
  - Quantize outputs to coarser levels
  - Ex: two vs. ten potential packet sizes
- What about performance?
  - Padding VoIP can introduce >42% overhead!!
  - Padding to 1500 bytes induces up to 150% overhead for HTTP!!



# **Implementation Problems**

- Performance almost always wins over security for consumer implementations
- Results in attacks that infer contents of encrypted network traffic
  - Login passwords for SSH  $\rightarrow$  packet timing
  - Web page identities for SSL/TLS  $\rightarrow$  packet size
  - Language and phrases for VoIP  $\rightarrow$  packet size



#### Overview

- Illustrate potential privacy problems introduced by network traffic features:
  - Uncovering Spoken Phrases in Encrypted
     Voice over IP Conversations
- Present efficient countermeasure by moving from information theoretic model:
  - Traffic Morphing: An Efficient Defense Against Statistical Traffic Analysis













Stream cipher used to efficiently encrypt packets









### **Attack Intuition**



#### **Attack Intuition**

#### **Length-preserving stream cipher keeps correlation**



#### **Attack Intuition**



Packet sizes carry information about the original audio

Original audio is made up of the building blocks of language called phonemes



#### **Attack Intuition**





#### **Phrase Spotting**

- How do we use the correlation between packet size and audio?
  - Build a statistical model for the sequence of packet sizes in phrase
  - Use the model to detect when a phrase of interest has occurred
- Known as phrase spotting

# REDJACK **Phrase Spotting** Packets:











# REDJACK **Phrase Spotting** Packets: Phrases: "spot me" "spot me" **Multiple** Sequence Insertion Alignment Deletion

# REDJACK **Phrase Spotting** Packets: Phrases: "spot me" "spot me" Aligned columns known as the consensus sequence for the phrase **Represented** with a **Profile Hidden Markov Model (HMM)**



- How do we train our phrase HMM without speaker-specific data?
- What if we do not have examples of the phrase we are looking for?

- Use concatenative synthesis!
  - Break phrase pronunciation into phonemes
  - Concatenate available phonemes together



# 1. Split phrase into words *the bike is red*





 Split phrase into words *the bike is red* 
 Break words into phonemes *dh ah b ay k ih z r eh d*





 Split phrase into words *the bike is red* 
 Break words into phonemes *dh ah b ay k ih z r eh d* 
 Replace phonemes with packet sizes





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- Split phrase into words the bike is red
   Break words into phonemes dh ah b ay k ih z r eh d
   Replace phonemes with packet sizes
- 4. Repeat to capture variations in speech



#### **Evaluation**

- TIMIT Dataset:
  - 462 training speakers, 168 testing speakers
  - 122 distinct phrases
  - Diverse set of dialects, pronunciations, etc.
- Use information retrieval metrics:
  - Recall = TP / (TP+FN)
  - Precision = TP / (TP+FP)



#### **Results**





#### **Results**

- Method is robust to variations in the audio
  - Noise
  - Gender
  - Dialect
  - Wide-band vs. Narrow-band
- Surprising amount of information leakage from encrypted VoIP traffic
  - More recent work produces transcripts



#### Countermeasures

- Padding can be effective for some types of information, but not others
- Can we do better?

Protocol	Leaked Info.	Block Size	Accuracy	Overhead
VoIP	Phrases	256-bit	0.4%	16.5%
	Language	512-bit	6.9%	<b>42.2</b> %
нттр	Web Pages	1500 bytes	35.9%	<b>140.8</b> %



# **Traffic Morphing**

- What if we could morph one class of traffic to look like another?
  - Ex: Farsi appears to be English
- Use convex optimization to stochastically change packet sizes to:
  - Minimize overhead
  - Ensure distribution "looks" like the target



#### **Source Distribution**

(Packet Sizes for Farsi)

$$X = \begin{bmatrix} x_1 \\ x_2 \\ \dots \\ x_n \end{bmatrix}$$

**Target Distribution** 

(Packet Sizes for English)

$$Y = \begin{bmatrix} y_1 \\ y_2 \\ \dots \\ y_n \end{bmatrix}$$





#### **Find morphing matrix A such that:**

$$\begin{bmatrix} y_1 \\ y_2 \\ \dots \\ y_n \end{bmatrix} = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \dots & \dots & \dots & \dots \\ a_{n1} & a_{n2} & \dots & a_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \dots \\ x_n \end{bmatrix}$$





Treat each column as a probability distribution for morphing source packet size s<sub>i</sub> to target

$$\begin{bmatrix} y_1 \\ y_2 \\ \dots \\ y_n \end{bmatrix} = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \dots & \dots & \dots & \dots \\ a_{n1} & a_{n2} & \dots & a_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \dots \\ x_n \end{bmatrix}$$



- Underspecified system of equations:
  - n<sup>2</sup> unknowns and 2n equations
  - Infinitely many solutions
- Choose one that minimizes the cost:

minimize 
$$f_0(A)$$
  
subject to  $\sum_{j=1}^n a_{ij} x_j = y_i, \quad \forall i \in [1, n]$   
 $\sum_{i=1}^n a_{ij} = 1, \quad \forall j \in [1, n]$   
 $a_{ij} \ge 0, \quad \forall i, j \in [1, n]$ 





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# **How to Morph Traffic**

- To morph source packet with size s<sub>i</sub>:
  - Find j<sup>th</sup> column of morphing matrix A
  - Sample from column according to distribution
  - Alter packet size to match sampled size s<sub>i</sub>

$y_1$		$a_{11}$	$a_{12}$	 $a_{1n}$	$\begin{bmatrix} x_1 \end{bmatrix}$
$y_2$	_	$a_{21}$	$a_{22}$	 $a_{2n}$	$x_2$
$y_n$		$a_{n1}$	$a_{n2}$	 $a_{nn}$	$x_n$



# **Real-world Challenges**

- Many issues arise in the real-world:
  - Reducing packet sizes?
  - Large sample spaces?
  - Slight changes in distribution?
  - Not enough packets?





# **VoIP Morphing Results**

- Apply morphing to VoIP language identification work of Wright et al.
  - Nearest-neighbor classifier
  - n-gram distributions of packet sizes
  - Can the classifier distinguish morphed from real?

Strategy	Overhead	Bigram Accuracy
No Padding	0.0%	71%
512-bit Padding	42.2%	50%
Morphing	15.4%	54%



# **VoIP Morphing Results**

- Morpher must be at least as "powerful" as the classifier to reduce accuracy
  - Trigram classifier beats bigram morphing
  - ...but we can make a trigram morpher...

Strategy	Bigram Accuracy	Trigram Accuracy	
No Padding	71%	76%	
512-bit Padding	50%	50%	
Morphing	54%	76%	



# **Open Questions**

- How do we determine security of morphing-like countermeasures?
  - No obvious proof since we are not working in an information theoretic model
- Is there a way to tell if a classifier can ever win?
  - Certain natural limitations and costs may prevent morphing in some scenarios
  - Ex: Cannot buffer packets indefinitely



#### Summary

- Most users assume contents are secure with well-known cryptographic protocols
- Performance considerations lead to information leakage from traffic features
- Explored security/performance trade-off:
  - Phrase spotting in encrypted VoIP
  - Countermeasures using traffic morphing

#### **Search Hidden Markov Model**



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# Search Hidden Markov Model



REDJACK







#### **Search Hidden Markov Model**

